# Lab 3

# **Interpreters and Types**

## **Objective**

- Understand visitors as a way to traverse a tree.
- Implement a typer and an interpreter as visitors.
- This is due on https://etudes.ens-lyon.fr (NO EMAIL PLEASE), before 2024-10-14 23:59. More instructions in section 3.6.

#### **EXERCISE** #1 ► Lab preparation

In the cap-lab24 directory: git pull will provide you all the necessary files for this lab in TP03 and MiniC. The latter folder will also be used for the next labs. ANTLR4 and pytest should be installed and working like in Lab 2, if not <sup>1</sup>:

```
python3 -m pip install pytest pytest-cov pytest-xdist
python3 -m pip install --upgrade coverage
```

## 3.1 Demo: Implicit tree walking using Visitors

## 3.1.1 Interpret (evaluate) arithmetic expressions with visitors

In the previous lab, we used an "attribute grammar" to evaluate arithmetic expressions during parsing. Today, we are going to let ANTLR build the syntax tree entirely, and then traverse this tree using the *Visitor* design pattern<sup>2</sup>. A visitor is a way to separate algorithms from the data structure they apply to.

For every possible type of node in your AST, a visitor will implement a method that will apply to nodes of this type.

## EXERCISE #2 ➤ Demo: arithmetic expression interpreter (TP03/arith-visitor/)

Observe and play with the Arit. q4 grammar and its PYTHON Visitor on myexample:

\$ make && make ex

Note that unlike the "attribute grammar" version that we used previously, the . g4 file does not contain Python code at all.

Have a look at the AritVisitor.py, which is automatically generated by ANTLR4: it provides an abstract visitor whose methods do nothing except a recursive call on children. Have a look at the MyAritVisitor.py file, observe how we override the methods to implement the interpreter, and try using print instructions to observe how the visitor actually works.

Also note the #blabla pragmas after each rules in the g4 file. They are here to provide ANTLR4 a name for each alternative in grammar rules. These names are used in the visitor classes, as method names that get called when the associated rule is found (eg. #foo will get visitFoo(ctx) to be called).

We depict the relationship between visitors' classes in Figure 3.1.

<sup>&</sup>lt;sup>1</sup>The second line is not always needed but may solve compatibility issues between versions of pytest-cov and coverage, yielding pytest-cov: Failed to setup subprocess coverage messages in some situations.

<sup>&</sup>lt;sup>2</sup>https://en.wikipedia.org/wiki/Visitor\_pattern

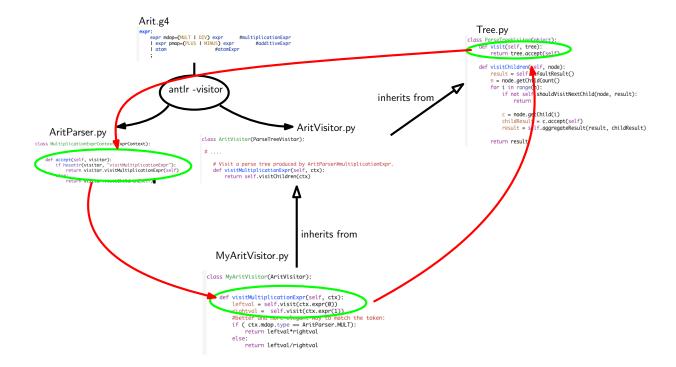


Figure 3.1: Visitor implementation Python/ANTLR4. ANTLR4 generates AritParser as well as AritVisitor. This AritVisitor inherits from the ParseTree visitor class (defined in Tree.py of the ANTLR4-Python library, use find to search for it). When visiting a grammar object, a call to visit calls the highest level visit, which itself calls the accept method of the Parser object of the good type (in AritParser) which finally calls your implementation of MyAritVisitor that match this particular type (here Multiplication). This process is depicted by the red cycle.

#### 3.1.2 Basic rules to write an ANTLR4 visitor

- For each alternative of each rule labeled #fooBar, write a method visitFooBar(self, ctx) where ctx is the corresponding parse subtree. Note the case change, #fooBar has lower-case f but visitFooBar has upper-case F.
- For each element elem in the right-hand side of the rule (lower-case non-terminal or upper-case terminal), you can access its value with ctx.elem(). For non-terminal (lower-case) elements, the corresponding value is a tree. For terminal (upper-case) elements, it is a token of type antlr4. Token. CommonToken. Tokens have in particular a field type whose value is AritParser. token-name.
- When the element appears several times, access the *n*-th instance using ctx.elem(*n*) (starting with 0). Note that when there's only one instance, you cannot use ctx.elem(0) but can only write ctx.elem().
- When the element is named like left=expr in the rule, access it with ctx.left (no parenthesis this time).
- Both trees and tokens have a getText() method returning the corresponding text in the source code.
- Recursive calls on sub-trees are written as self.visit(subtree).
- what you can write in Python code is dictated by the .g4 file.

Example: when a ANTLR4 rule contains an operator alternative such as:

```
| expr addop=(PLUS | MINUS) expr #additiveExpr
```

you can use the following code in your implementation of visitAdditiveExpr to match the operator:

```
def visitAdditiveExpr(self, ctx):
   leftval = self.visit(ctx.expr(0))
   rightval = self.visit(ctx.expr(1))
   if ctx.addop.type == AritParser.PLUS:
      return leftval + rightval
   else:
      return leftval -rightval
```

Note that we wrote PLUS and MINUS in the same rule to have the same level of precedence, and avoid the issues we had in lab 2.

## 3.2 Up to you: first visitors

## **EXERCISE** #3 ► Trees (should be quickly done!)

Consider the following grammar:

This grammar represents "scheme-like trees", for instance node (42 12 1515 17) is the tree with root 42 and three children 12, 1515, 17.

- 1. We give you the grammar in the folder tree/. Copy and adapt previous files to get it operational.
- 2. Implement a visitor that decides whether a syntactically correct file is a binary tree. Your main file should contain:

```
tree = parser.int_tree_top()
visitor = MyTreeVisitor()
is_binary_tree: bool = visitor.visit(tree)
print("Is_it_a_binary_tree_?" + str(is_binary_tree))
```

## 3.2.1 Application to MiniC Language

The objective is now to use visitors, to type and interpret MiniC programs, whose syntax is depicted in Figure 3.2. Classically, we should do typing first and the interpretation afterwards, but in this lab we will implement the interpretation first (assuming the program is well-typed).

```
grammar MiniC;
prog: function* EOF #progRule;
// For now, we don't have "real" functions, just the main() function
// that is the main program, with a hardcoded profile and final // 'return 0' (actually a 'return INT' because we don't have a ZERO
// lexical token).
function: INTTYPE ID OPAR CPAR OBRACE vardecl_1 block
    RETURN INT SCOL CBRACE #funcDef;
vardecl_l: vardecl* #varDeclList;
vardecl: typee id_l SCOL #varDecl;
id_1: ID #idListBase
    | ID COM id_l #idList
block: stat* #statList;
stat: assignment SCOL
    | if_stat
      while_stat
    | print_stat
assignment: ID ASSIGN expr #assignStat;
if_stat: IF OPAR expr CPAR then_block=stat_block
        (ELSE else_block=stat_block)? #ifStat;
stat_block: OBRACE block CBRACE
         stat
while_stat: WHILE OPAR expr CPAR body=stat_block #whileStat;
print_stat
    : PRINTLN_INT OPAR expr CPAR SCOL #printlnintStat
     PRINTLN_FLOAT OPAR expr CPAR SCOL #printlnfloatStat
     PRINTLN_BOOL OPAR expr CPAR SCOL #printlnboolStat
    | PRINTLN_STRING OPAR expr CPAR SCOL #printlnstringStat
```

Figure 3.2: MiniC syntax. We omitted here the subgrammar for expressions

### EXERCISE #4 ► Be prepared!

In the directory MiniC/ (outside TP03/), you will find:

- The MiniC grammar (MiniC.g4). Run make to run ANTLR on it and generate the corresponding Python files.
- Our "main" program (MiniCC.py) which does the parsing of the input file, then launches the Typing visitor, and if the file is well typed, launches the Interpreter visitor. In this lab it supports four modes:
  - python3 MiniCC.py --mode parse <file> checks whether the given file is syntactically valid MiniC code.
  - python3 MiniCC.py --mode typecheck <file> parses the given file and typechecks it.
  - python3 MiniCC.py --mode eval <file> parses, typechecks, and interprets the given program.
  - python3 MiniCC.py --mode eval --disable-typecheck <file> parses and interprets the given program, but does not typecheck it. This will be useful before you complete the typechecker implementation.

Try it on some provided examples (e.g. in TP03/tests/provided/examples-types/), see what happens for well-typed and ill-typed programs (usually named bad\_\*.c).

- Two visitors to be completed: TP03/MiniCTypingVisitor.py and TP03/MiniCInterpretVisitor.py.
- Some test cases (TP03/tests), and a test infrastructure.

## 3.3 An interpreter for the MiniC-language

### 3.3.1 Informal Specifications of the MiniC Language Semantics

MiniC is a small imperative language inspired from C, with more restrictive typing and semantic rules. Some constructs have an undefined behavior in C and well defined semantics in MiniC:

- Variables that are not explicitly initialized in the program are automatically initialized:
  - to 0 for int,
  - to 0.0 for float,
  - to false for bool,
  - to the empty string "" for string.
- Divisions and modulo by 0 must print the message "Division by 0" and stop program execution with status 1 (use raise MiniCRuntimeError("Division by 0") to achieve this in the interpreter).
- Conventions for division and modulo are the same as in C: division is truncated toward zero, and modulo is such that (a/b) \* b + a%b = a.

$$4/3 = 1$$
  $4\%3 = 1$   $(-4)/3 = -1$   $(-4)\%3 = -1$   $4/(-3) = -1$   $4\%(-3) = 1$   $(-4)/(-3) = 1$   $(-4)/(-3) = -1$ 

#### 3.3.2 Implementation of the Interpreter

The semantics of the MiniC language (how to evaluate a given MiniC program) is defined by induction on the syntax. You already saw how to evaluate a given expression, this is depicted in Figure 3.3.

#### **EXERCISE** #5 ► **Interpreter rules (on paper)**

First fill the empty cells in Figure 3.4, then ask your teaching assistant to correct them.

#### EXERCISE #6 ► Interpreter

Now you have to implement the interpreter of the MiniC language. We give you the structure of the code and the implementation for numerical expressions and boolean expressions (except modulo!). You can reason in terms of "well-typed programs", since badly typed programs should have been rejected earlier.

Run the command:

literal constant c	return int(c) or float(c)
variable name x	find value of x in dictionary and return it
$e_1+e_2$	v1 <- e1.visit() v2 <- e2.visit() return v1+v2
true	return true
$e_1 < e_2$	return e1.visit() <e2.visit()< td=""></e2.visit()<>

Figure 3.3: Interpretation (Evaluation) of expressions

x := e	<pre>v &lt;- e.visit() store(x,v) #update the value in dict</pre>
println_int(e)	<pre>v &lt;- e.visit() print(v) # python's print</pre>
S1; S2	s1.visit() s2.visit()
if b then S1 else S2	
while $b$ do $S$ done	

Figure 3.4: Interpretation for Statements (pseudo-code)

make
python3 MiniCC.py --mode eval --disable-typecheck TP03/tests/provided/examples/test\_print\_int.c

and the interpreter will be run on test\_print\_int.c. On the particular example test\_print\_int.c observe how integer values are printed.

You still have to implement (in MiniCInterpretVisitor.py):

- 1. The modulo version of Multiplicative expressions (for the C language semantics of modulo).
- 2. Variable declarations (varDecl) and variable use (idAtom): your interpreter should use a table (*dict* in PYTHON) to store variable definitions and check if variables are correctly defined and initialized. **Do not forget to initialize dict with the initial values (0, 0.0, False or "" depending on the type) for all variable declarations.**
- 3. Statements: assignments, conditional blocks, tests, loops.

**Error codes** The exit code of the interpreter should be:

- 1 in case of runtime error (e.g. division by 0, absence of main function)
- 2 in case of typing error
- 3 in case of syntax error
- 4 in case of internal error (i.e. error that should never happen except during debugging)
- 5 in case of unsupported construct (should not be used in lab3, but you will need it for strings and floats during code generation)
- And obviously, 0 if the program is typechecked and executed without errors.

The file MiniCC.py in the skeleton already does this for you if you raise the right exception (see Lib/Errors.py). You need to use these codes as test annotations in programs raising errors:

- Programs raising a runtime error should be annotated with // EXECCODE 1
- Programs rejected by the interpreter before execution should be annotated with // EXITCODE *n*, with *n* being 2, 3, 4 or 5 as documented above.

The distinction between EXECCODE and EXITCODE seems subtle for an interpreter, but will be more obvious for a compiler, where EXITCODE will refer to the exit code of the *compiler*, and EXECCODE to the exit code of the *program's execution*.

#### **EXERCISE** #7 ► Automated tests

We provide a script to automatically test your code. As at this point, you do not have your typechecker, modify the variable DISABLE\_TYPECHECK to True in test\_interpreter.py. You will have to put it back to False once you begin working on MiniCTypingVisitor.py.

Test with make test and write an appropriate test-suite. If you get an error about the --cov argument, you didn't properly install pytest-cov. You must provide your own tests: they will be graded depending on their quality. The only outputs are the one from the println\_\* function or the following error messages: "m has no value yet!" (or possibly "Undefined variable m", but this error should never happen if your typechecker did its job properly) where m is the name of the variable. In case the program has no main function, the typechecker accepts the program, but it cannot be executed, hence the interpreter raises a "No main function in file" runtime error. (Note that error messages raised from the typechecker have stricter formatting requirements, see below.)

**Test Infrastructure** Tests work mostly as in the previous lab. By default, the testsuite is ran on all .c files in the TP03/tests/directory. You may restrict to a set of files using make test FILTER=.... FILTER is either a single file or an extended wildcard like TP03/tests/provided/\*\*/\*.c (\*\* matches any directory hierarchy).

Source files should contain // EXPECTED and // EXITCODE n pragmas to specify the expected behavior of the compiler. They are special comments (the // is needed to keep compatibility with C, only the testsuite considers them as special). The EXITCODE corresponds to the exit codes described in Section 3.3.2.

For instance, if you fail test\_print\_int.c because you printed 43 instead of 42, using the command make test FILTER=TP03/tests/provided/examples/test\_print\_int.c you will get this error:

```
_TestInterpret.test_eval[TP03/tests/provided/examples/test_print_int.c] _____
self = <test_interpreter.TestInterpret object at 0x7f05d6f86a40>,
filename = 'TP03/tests/provided/examples/test_print_int.c'
    @pytest.mark.parametrize('filename', ALL_FILES)
    def test_eval(self, filename):
        cat(filename) # For diagnosis
        expect = self.get_expect(filename)
        eval = self.evaluate(filename)
        if expect:
            self.assert_equal(eval, expect)
test interpreter.pv:48:
self = <test_interpreter.TestInterpret object at 0x7f05d6f86a40>,
actual = testinfo(exitcode=0, execcode=0, output='43\n', linkargs=[], skip_test_expected=False)
expected = testinfo(exitcode=0, execcode=0, output='42\n', linkargs=[], skip_test_expected=False)
    def assert_equal(self, actual, expected):
        if expected.output is not None and actual.output is not None:
            assert actual.output == expected.output, \
>
                "Output of the program is incorrect."
F.
            AssertionError: Output of the program is incorrect.
            assert '43\n' == '42\n'
Ε
Ε
              - 42
E.
              + 43
test_interpreter.py:35: AssertionError
   And if you did not print anything at all when 42 was expected, the last lines would be this instead:
    def assert_equal(self, actual, expected):
        if expected.output is not None and actual.output is not None:
            assert actual.output == expected.output, \
>
                "Output of the program is incorrect."
            AssertionError: Output of the program is incorrect.
F
            assert '' == '42\n'
F
              - 42
test_interpreter.py:35: AssertionError
```

## 3.4 A type-checker for the MiniC language

#### 3.4.1 Informal Typing Specification for the MiniC Language

MiniC is a subset of C with stricter rules, and predefined aliases:

```
typedef char * string;
typedef int bool;
static const int true = 1;
static const int false = 0;
```

The informal typing rules for the MiniC language are:

- Variables must be declared before being used, and can be declared only once;
- Binary operations (+, -, \*, ==, !=, <=, &&, | |, ...) require both arguments to be of the same type (e.g. 1 + 2.0 is rejected);
- Boolean and integers are incompatible types (e.g. while (1) is rejected);
- Binary arithmetic operators return the same type as their operands (e.g. 2. + 3. is a float, 1 / 2 is the integer division);
- Modulo (%) is accepted only on integers, not floats.

- + is accepted on string (it is the concatenation operator), no other arithmetic operator is allowed for string;
- Comparison operators (==, <=, ...) and logic operators (&&, | |) return a Boolean;
- == and != accept any type as operands;
- Other comparison operators (<, >=, ...) accept int and float operands only.

The expected errors of the typechecker are the following:

- "In function f: Line l col c: type mismatch for e: t1 and t2" for assignments and comparison (equality operands only), if the two arguments have different types;
- "In function f: Line l col c: invalid type for MESSAGE: t (and t')" for typing error, with MESSAGE explicit enough. For example: "In function main: Line 8 col 6: invalid type for multiplicative operands: integer and string";
- "In function f: Line l col c: MESSAGE" for errors that are not purely typing, e.g. undeclared variable or double declared variables. For example: "In function main: Line 5 col 2: Variable x already declared".

The name f is the current function, for the moment it should be 'main' but we may add functions later. Some of the tests provided, mainly of the form bad\_...c, give examples of expected errors.

As before, we explicitly ask you to write new test cases, and make your error messages as explicit as possible.

## 3.4.2 Implementation of the Typechecker

#### EXERCISE #8 ► Typing

Write typing rules for expressions (on paper). Then, implement a type checker for the MiniC language<sup>3</sup> (as a standalone visitor MiniCTypingVisitor). We provide you with a (basic) class for basic types and the environment initialization with the declared types. The methods \_raise, \_raiseNonType and \_assertSameType allow you to add informative exception handlers. The provided test files must guide you when the implementation cannot be directly derived from the typing rules. Testing is the same as for the interpreter, except that you do not have to put --disable-typecheck on your individual test.

Do not forget to modify the variable DISABLE\_TYPECHECK back to False in test\_interpreter.py, otherwise the command make test will not test your MinicTypingVisitor.py!

Also, do not forget to intensively use new test files.

## 3.5 Language extensions

In this section, the instructions are all the same: for each new extension, implement the syntax, give new semantic rules (on paper), give new interpretation rules (code), new typing rules, relevant test cases, adapt the test infrastructure, . . . .

The maximum grade (20/20) correspond to a clear and documented code without any flaw, implementing at least one of the following extensions, and with a test suite of quality.

#### **EXERCISE** #9 ► Fortran-like for loops

Implement typing and interpretation for loops that look like the following example (static loop bounds, optional constant stride):

k=42; for i=k to k+1515 by 2 { .... } Informal typing and semantics:

- The loop counter must be declared explicitly as int type before the loop;
- for i = a to b is an empty loop if b is strictly smaller than a (except with negative stride);
- Stride can be any integer value. When null, the loop is infinite.

 $<sup>^3</sup>$ We do not ask for a decorated AST, only type checking.

- Assigning the loop index within the loop is allowed, and when this happens the value assigned does not impact the next loop iterations (like Python's for i in range(...): loop).
- Loop bounds are evaluated when entering the loop, and not re-evaluated afterwards.

### **EXERCISE** #10 ► **C-like for loops**

Extend the language with C-like for loops, with initial assignment, loop condition and increment assignment all optional. Example are the followings:

```
for (i=1;i<4;i=i+1) { .... }
for (;j<4.0;j=j+1.5) { .... }
```

## 3.6 Final delivery

We recall that your work is **personal** and code copy (including tests) is **strictly forbidden**.

#### EXERCISE #11 ► Archive

The interpreter and the typer (working together) are due on the course's webpage

```
https://etudes.ens-lyon.fr/
```

Type make tar in the MiniC folder to obtain the archive MYNAME.tgz to send (change your name in the Makefile before!). You have a (minimal) README-interpreter.md to fill with your name, the functionality of the code, how to use it, your design choices if any, the chosen extensions, and known bugs. Your archive must also contain your tests (TESTS!) in the TP03/tests/students folder. We expect unit tests (small files that test just one feature, or the interaction between a few ones, not everything at once in a huge test file) with clear and explicit names (typically test\_str\_assign.c and not test04.c). The command make test must work with your implementation; if some of the tests you have fail, please report the corresponding bugs in your readme. When there are some error messages given in examples and tests we provide, you are expected to write your code to produce exactly these messages.

You will be graded on 20 points, as follows:

- 10 points on the correctness of your code (do you pass all tests we could think of?),
- 5 points on the coverage of your tests (when using make tests with the provided tests and yours, are there lines of **our** code which are never executed?),
  - → you can check your coverage by going to MiniC/htmlcov/index.html; of course, you should cover only the code you are supposed to write, and not for instance the code in MiniCC.py,
- 5 points on the correctness and coverage of the chosen extension (if any), the quality of your code and tests, your readme, your archive,...,
- if you deposit your work late, you will lose 1 point per hour of lateness,
- in case of plagiarism, if *n* students have the same code, then the grade of each student will be divided by *n*.